

Recall from last time.

Thm (recurrence state thm).

State i is recurrent $\Leftrightarrow$

$$\sum_{n=1}^{+\infty} P_{ii}^{(n)} = +\infty.$$

(w.p.1, MC starting from i will return to i infinitely often)

$P_{ii}^{(n)}$ = n -step transition.

eg. application, multidimensional RW.

$$\sum_{n=1}^{+\infty} P_{ii}^{(n)} = \sum_{n=1}^{+\infty} \left(\frac{c}{\sqrt{n}}\right)^d \begin{cases} < +\infty & d \geq 3 \\ = +\infty & d = 1, 2. \end{cases}$$

Proof of the theorem.

Detour: Borel-Cantelli Lemma.

General tool for reasoning infinite seq of events
(finitely or infinitely often)

Let $(E_n)_{n=1}^{+\infty}$ be a sequence of events,

If $\sum_{n=1}^{+\infty} P(E_n) < +\infty$

then

$$P\left(\left(E_n\right)_{n=1}^{+\infty} \text{ happens only finitely often}\right) = 1.$$

Note: converse is false in general.

Application to recurrence state thm:

We let $E_n = \{X_n = i\}$ (visit back to i at time n).

$$\text{If } \sum_{n=1}^{+\infty} P(E_n) = \sum_{n=1}^{+\infty} p_{ii}^{(n)} < +\infty$$

then by BC lemma. $\left(E_n\right)_{n=1}^{+\infty}$ happens only finitely many times, w.p.1

i.e. we only visit back to i finitely

many times, w.p.1.

This implies transience. and proves " $\Rightarrow$ " of thm.

Back to the detour.

Proof of BC lemma.

$$\sum_{n=1}^{+\infty} P(E_n) = \underbrace{E}_{\text{Fubini-Tonelli!}} \left[\sum_{n=1}^{+\infty} 1_{E_n} \right] < +\infty \quad (\text{by assumption}).$$

↑
of times event happens.

So the r.v. $\sum_{n=1}^{\infty} 1_{E_n}$ has finite expectation.

and thus finite a.s.

From the proof, it's clear that converse cannot be guaranteed — $\exists$ real-valued r.v. w/ infinite $E[\dots]$.

However, for MC's. $N(i) = \sum_{n=1}^{\infty} 1_{\{X_n=i\}}$

is either infinite w.p.1 or

finite & geometrically distributed.

Using this observation, we can prove " $\Leftarrow$ " of recurrence state thm.

$$\sum_{n=1}^{\infty} p_{ii}^{(n)} \underset{(\text{Fubini-Tonelli})}{=} E_i \left[\sum_{n=1}^{\infty} 1_{X_n=i} \right]$$

$$= E_i [N(i)].$$

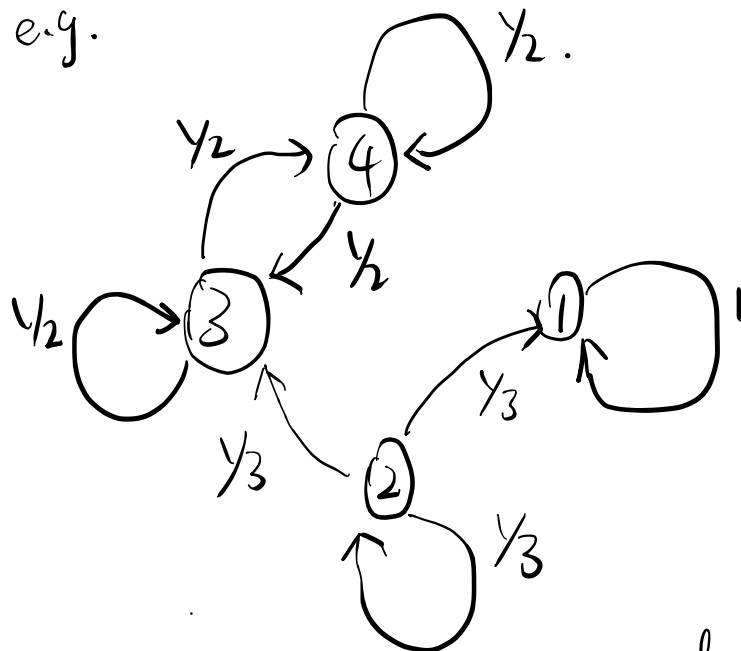
If i is transient, then $f_{ii} < 1$

$$E_i [N(i)] = \frac{f_{ii}}{1 - f_{ii}} < +\infty.$$

This proves " $\Leftarrow$ " of thm.

Next question: how to compute f_{ij} 's systematically?

Motivating e.g.



$$f_{11} = 1. \quad f_{ij} = 0 \quad \text{for } j \in \{2, 3, 4\}.$$

$$f_{22} = \frac{1}{3}.$$

$$f_{21} = \frac{1}{3} + \frac{1}{3} \times \frac{1}{3} + \left(\frac{1}{3}\right)^3 + \dots$$

$$= \frac{\frac{1}{3}}{1 - \frac{1}{3}} = \frac{1}{2}.$$

$$\text{Similarly. } f_{23} = f_{24} = \frac{1}{2}$$

In general, we use f -expansion.

i.e. solving out f_{ij} 's using linear systems.

$$f_{ij} = \mathbb{P}_i(\text{visit } j).$$

(Expand based on the first step)

$$= \sum_{k \in S} \mathbb{P}_i(X_1 = k) \cdot \mathbb{P}_k(\text{visit } j \text{ starting from time } 0).$$

($\neq \sum_{k \in S} P_{ik} f_{kj}$)

$$= \sum_{\substack{k \in S \\ k \neq j}} P_{ik} f_{kj} + P_{ij}$$

If $|S| < \infty$, f -expansion gives $|S|^2$ equations for $|S|^2$ variables.

We can solve out these variables.

A more interesting example.

— Gambler's ruin.

Suppose a gambler plays coin-tossing games.

- Initially, $a \in \mathbb{N}_+$ units of money.
- For each round, bet 1 unit

with winning prob $p \in (0, 1)$.

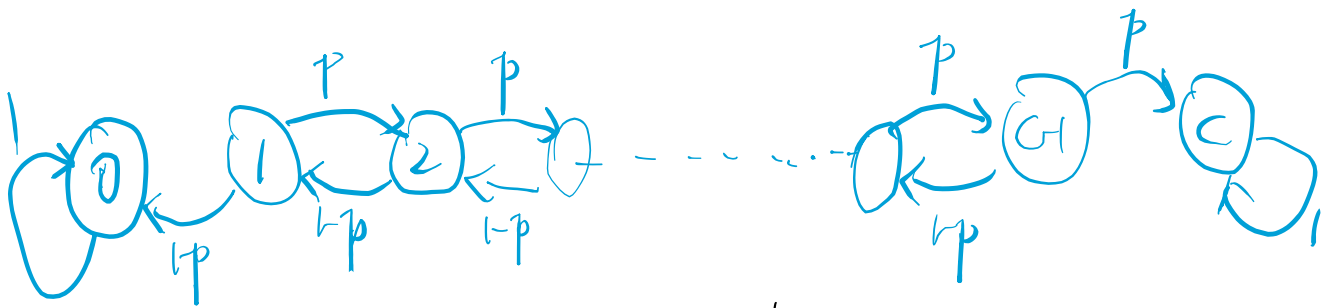
- Stop when $\begin{cases} \text{money} = 0 \\ \text{money} = c > a. \end{cases}$

In MC language, let X_n be the amount of money in n -th round.

$$X_0 = a. \quad X_{n+1} = X_n + \epsilon_{n+1} \text{ when } X_n \in \{1, 2, \dots, c\}$$

$$\text{where } \epsilon_{n+1} = \begin{cases} +1 & \text{w.p. } p \\ -1 & \text{w.p. } (1-p) \end{cases} \text{ indep.}$$

$$X_{n+1} = X_n \text{ when } X_n \in \{0, c\}.$$



Interested in ruin probability

$$\mathbb{P}_a(\text{visit } 0) = f_{a0}.$$

Computation of f_{a0} :

$$f_{00} = 1 \quad f_{c0} = 0$$

By f -expansion.

$$f_{a0} = p f_{(a+1)0} + (1-p) f_{(a-1)0}$$

(when $a \geq 2$)

For $a=1$

$$f_{10} = p f_{20} + (1-p) f_{00}$$

($= p f_{20} + (1-p) f_{00}$ since $f_{00}=1$)

So $f_{a0} = p f_{(a+1)0} + (1-p) f_{(a-1)0}$ for $a \in \{1, 2, \dots, c\}$

Solving the f -expansion eqs.

Case I: $p = \frac{1}{2}$.

$$f_{a0} = \frac{1}{2} (f_{(a+1)0} + f_{(a-1)0}),$$

So $(f_{a0})_{a=0}^c$ is an arithmetic sequence, i.e.

$$f_{c0} - f_{(c-1)0} = f_{(a+1)0} - f_{a0} = f_{a0} - f_{(a-1)0} = \dots = f_{10} - f_{00}$$

$\parallel$
 $-\frac{1}{c}$

$$f_{00}=1, f_{c0}=0$$

and therefore,

$$f_{a0} = 1 - \frac{a}{c}.$$

Similarly, $f_{ac} = \frac{a}{c}.$

Remark: We cannot make money from it
(in expectation).

$$E[X_{\text{final}}] = c \cdot \frac{a}{c} + 0 \cdot \left(1 - \frac{a}{c}\right) = a.$$

(Indeed, this is not a coincidence).
C.f. martingale part.

Case II: $p \neq \frac{1}{2}.$

$$f_{a0} = p f_{(a+1)0} + (1-p) f_{(a-1)0}$$

$$p f_{(a+1)0} - p f_{a0} + (1-p) f_{(a-1)0} - (1-p) f_{a0} = 0.$$

So

$$f_{(a+1)0} - f_{a0} = \frac{1-p}{p} (f_{a0} - f_{(a-1)0}).$$

The increments form a geometric seq.

$$f_{(a+1)0} - f_{a0} = \left(\frac{1-p}{p}\right)^a (f_{10} - f_{00})$$

Summing up over a ,

$$\begin{aligned} -1 = f_{C0} - f_{00} &= \sum_{a=0}^{C-1} (f_{(a+1)0} - f_{a0}) \\ &= (f_{10} - f_{00}) \cdot \left[\sum_{a=0}^{C-1} \left(\frac{1-p}{p}\right)^a \right]. \end{aligned}$$

Following some calculation,

$$f_{a0} = \frac{\left(\frac{1-p}{p}\right)^C - \left(\frac{1-p}{p}\right)^a}{\left(\frac{1-p}{p}\right)^C - 1} \quad \left(p \neq \frac{1}{2}\right)$$

Remark: when $p \neq \frac{1}{2}$.

ruin prob is $\left\{ \begin{array}{l} \text{exponentially small when } p > \frac{1}{2}. \\ \text{exponentially close to 1 when } p < \frac{1}{2}. \end{array} \right.$

Relation between recurrence/transience of different states.

Recall notation:

$i \rightarrow j$ means j is reachable from i , $f_{ij} > 0$.

$i \leftrightarrow j$ means $i \rightarrow j$ and $j \rightarrow i$.

Def. We call an MC irreducible when

$i \leftrightarrow j$ for $i, j \in S$.

(Implicitly, an MC not satisfying this
can be "reduced" into smaller ones).

Fact. If $i \leftrightarrow j$, then " i is recurrent"



" j is recurrent".

Corollary of the fact: "class theorem"

If a MC is irreducible, then one of the following is true:

- All states are recurrent. ("recurrent MC")
- All states are transient ("transient MC").

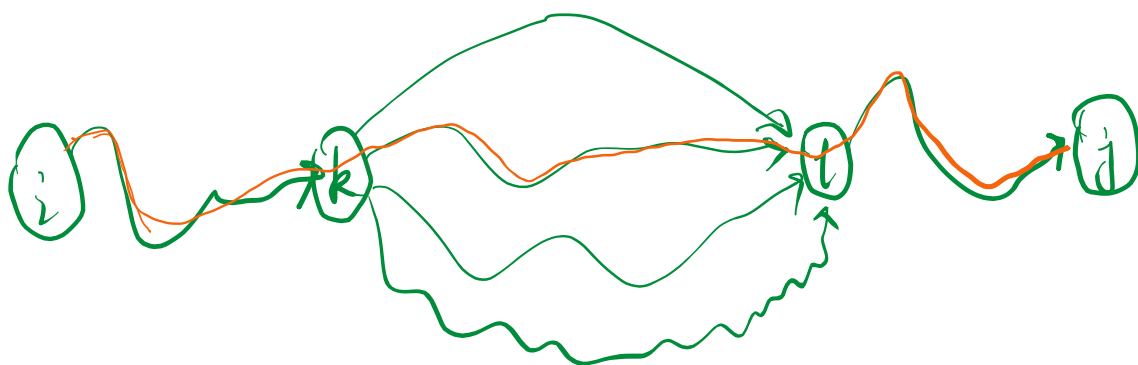
Proof of the fact:

We use the "sum lemma".

Lemma: If $i, j, k, l \in S$ (allow repeating).

Suppose $i \rightarrow k$ and $l \rightarrow j$

If $\sum_{n=1}^{+\infty} p_{kl}^{(n)} = +\infty$ then $\sum_{n=1}^{+\infty} p_{ij}^{(n)} = +\infty$.



Proof: Since $l \rightarrow k$, $\exists m > 0$, $p_{lk}^{(m)} > 0$

$l \rightarrow j$, $\exists r > 0$, $p_{lj}^{(r)} > 0$.

Then (by Kolmogorov - Chapman equation)

$$p_{ij}^{(n+m+r)} \geq p_{ik}^{(m)} \cdot p_{kl}^{(n)} \cdot p_{lj}^{(r)}.$$

um them together

$$\sum_{n=1}^{+\infty} p_{ij}^{(n)} \geq \underbrace{p_{ik}^{(m)}}_{>0} \cdot \underbrace{p_{lj}^{(r)}}_{>0} \cdot \underbrace{\sum_{n=1}^{+\infty} p_{kl}^{(n)}}_{=+\infty} = +\infty.$$

This is sum lemma

I dk to prove of recurrence equivalence:

Apply sum lemma w/ $j=i, l=k$.

$$i \rightarrow k, l \rightarrow j \iff i \leftrightarrow l.$$

$$l \text{ is recurrent} \iff \sum_{n=1}^{+\infty} p_{ll}^{(n)} = +\infty.$$

$$\text{sum lemma implies} \quad \sum_{n=1}^{+\infty} p_{ll}^{(n)} = +\infty.$$

which proves desired fact.

Next question.

What do we know about $\sum_{n=1}^{+\infty} p_{ij}^{(n)}$? (for $i \neq j$).

Fact. If MC is $\begin{cases} \text{Recurrent,} \\ \text{Transient,} \end{cases}$ $\sum_{n=1}^{+\infty} p_{ij}^{(n)} = +\infty$ $\sum_{n=1}^{+\infty} p_{ij}^{(n)} < +\infty$.

(This is if-and-only-if: any if pair
is equivalent to MC recurrence/transience)

Proof. $\sum_{n=1}^{+\infty} P_{ij}^{(n)} = \mathbb{E}_i[N(j)]$.

• Recurrent, $\exists m, P_{ij}^{(m)} > 0$ by irreducibility

$$\mathbb{E}_i[N(j)] \geq P_{ij}^{(m)} \cdot \underbrace{\mathbb{E}_j[N(j)]}_{= +\infty \text{ by recurrence}} = +\infty.$$

• Transient.

$$\begin{aligned} \mathbb{E}_i[N(j)] &= f_{ij} \cdot (\underbrace{\mathbb{E}_j[N(j)]}_{= \frac{f_{ji}}{1-f_{jj}}}) + 1 \\ &= \frac{f_{ij}}{1-f_{jj}} < +\infty. \end{aligned}$$

e.g. When $|S| < +\infty$, and irreducible,
then the MC is recurrent

Proof. For any fixed $i \in S$.

$$\underbrace{\sum_{j \in S} \left(\sum_{n=1}^{+\infty} P_{ij}^{(n)} \right)}_{\text{sum of finitely many terms.}} = \sum_{n=1}^{+\infty} \left(\sum_{j \in S} P_{ij}^{(n)} \right) = \sum_{n=1}^{+\infty} 1 = +\infty.$$

$\exists j_0$, s.t.

$$\sum_{n=1}^{+\infty} P_{ij_0}^{(n)} = +\infty.$$

and by irreducibility, this implies MC is recurrent.

Another useful fact: "f-lemma".

Motivation: $\exists$ characterization for recurrence/transience

$$\left\{ \begin{array}{l} \sum_{n=1}^{+\infty} P_{ij}^{(n)} \\ N(j) \\ f_{ij} ? \end{array} \right\} \text{ ————— equivalent}$$

the same

f Lemma. If $j \rightarrow i$, $f_{jj} = 1$

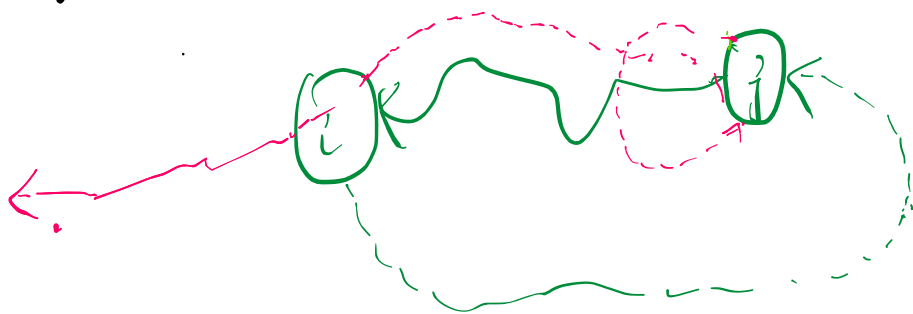
then $f_{ij} = 1$.

(Compute off-diagonal f for recurrent states)

Direct corollary:

For recurrent MC, $f_{ij} = 1$ for $i, j \in S$.

Proof of f -Lemma.



technical tool: "hit lemma".

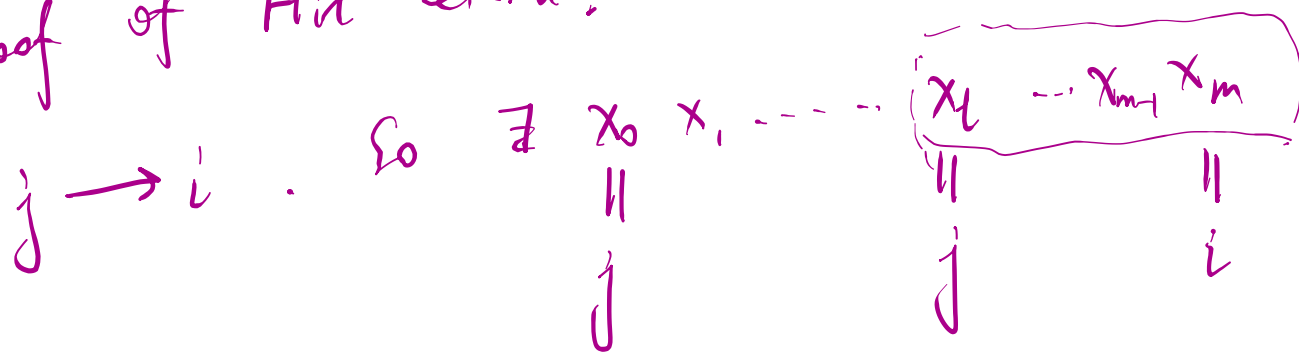
Define the event

$H_{ij} = \{ \text{MC hits } i \text{ before returning to } j \}.$

$f \ j \rightarrow i$, then $P_j(H_{ij}) > 0.$

(This lemma removes the trouble of " $\rightarrow$ " path).

Proof of "Hit Lemma".



We have

$$P_{x_0 x_1} P_{x_1 x_2} \dots P_{x_{m-1} x_m} > 0.$$

This implies

$$P_{x_1 x_{t+1}} P_{x_{t+1} x_{t+2}} \dots P_{x_{m-1} x_m} > 0.$$

Let t be the last time j appears
in this sequence.

$$P_j(H_j) \geq P_{x_1 x_{t+1}} \dots P_{x_{m-1} x_m} > 0$$

which proves hit lemma.

Proof of f -lemma using hit lemma.

$$0 = P_j(\text{never return to } j)$$

$$\geq \underbrace{P_j(H_j)}_{>0} \cdot \boxed{P_i(\text{never visit } j).}$$

must be 0.

So $f_{ij} = 1$, which proves f -lemma.